

# Interior-Point Methods

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## Before we go...

- Interior-point methods play an indispensable role in convex optimization.
- Modern LP/SOCP/SDP solvers, such as SeDuMi, SDPT3, and DSDP, are interior-point methods.

- What we will do:

Provide intuitive insights into the ideas that led to this beautiful technique.

- What we will not go through:

Convergence/complexity analysis, coverage of all interior-point algorithms (there are numerous!), complete derivations, implementation details, . . .

## Trapping Points within the Feasible Set

- Convex problem in standard form

$$\begin{aligned} \min \quad & f_0(x) \\ \text{s.t.} \quad & Ax = b, \quad f_i(x) \leq 0, \quad i = 1, \dots, m \end{aligned} \tag{†}$$

where  $f_0, \dots, f_m$  are assumed to be convex & twice differentiable.

- The inequality constraints can be made implicit by rewriting (†) as

$$\begin{aligned} \min \quad & f_0(x) + \sum_{i=1}^m I_-(f_i(x)) \\ \text{s.t.} \quad & Ax = b \end{aligned}$$

where  $I_-(u) = 0$  for  $u \leq 0$ ,  $I_-(u) = \infty$  otherwise.

- But  $I_-$  is not differentiable.
- **The basic idea:** approximate  $I_-$  by some differentiable function.

# Logarithmic Barrier Function

- Approximate  $I_-$  by

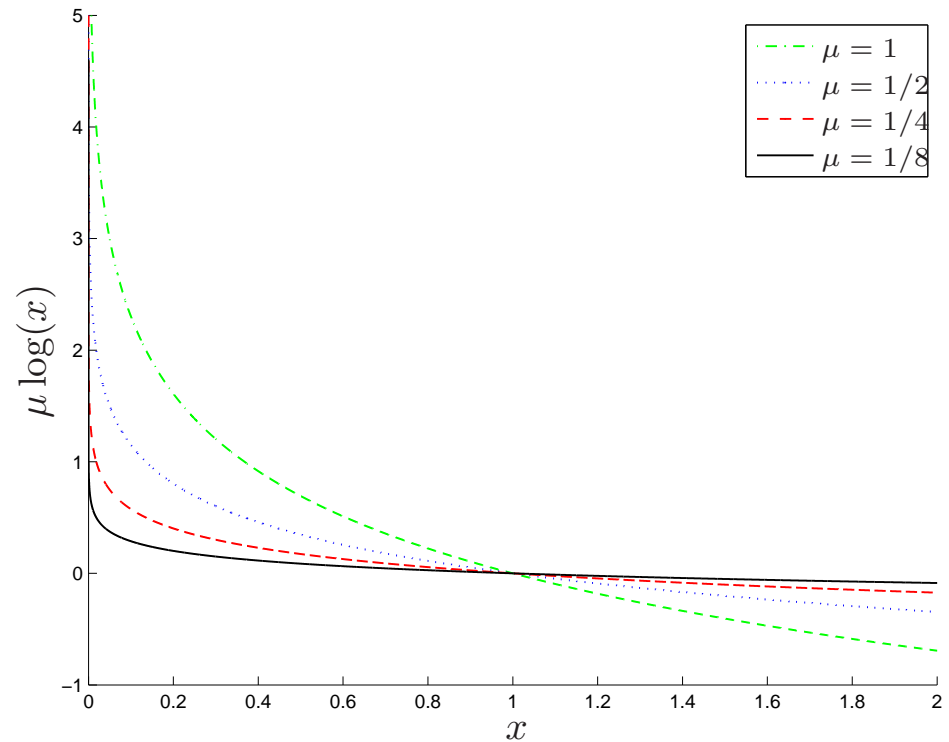
$$\hat{I}_-(u) = -\mu \log(-u), \quad \text{dom} \hat{I}_- = \{x \in \mathbf{R} \mid x < 0\}$$

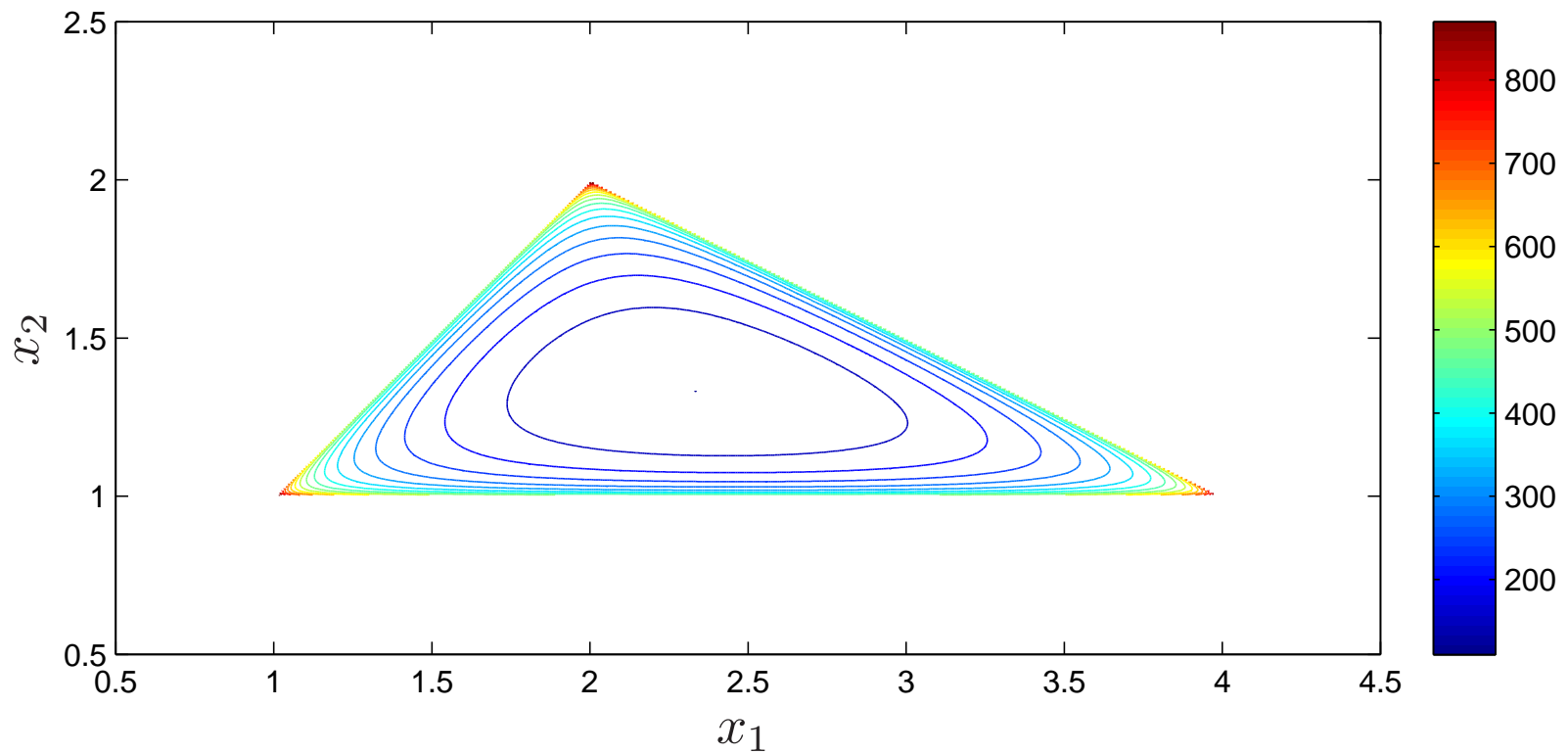
where  $\mu > 0$  is a parameter that controls the accuracy of the approx.

**Example:** a single inequality  
 $x \geq 0$ .

The corresponding approx. is

$$-\mu \log(x)$$





**Example:** A set of linear inequalities  $b_i - a_i^T x \leq 0$ ,  $i = 1, \dots, m$ .

The corresponding approx. is

$$-\mu \sum_{i=1}^m \log(a_i^T x - b_i)$$

- An approximation of the original problem

$$\begin{aligned} \min \quad & f_0(x) + \mu\phi(x) \\ \text{s.t.} \quad & Ax = b \end{aligned} \tag{†}$$

with  $\phi(x) = -\sum_{i=1}^m \log(-f_i(x))$ ,  $\text{dom}\phi = \{x \mid f_i(x) < 0, i = 1, \dots, m\}$ .

- $\phi$  is called the **logarithmic barrier function**. Some nice properties:
  - $\phi$  is convex (by composition).
  - $\phi$  is twice differentiable:

$$\begin{aligned} \nabla\phi(x) &= \sum_{i=1}^m \frac{1}{-f_i(x)} \nabla f_i(x) \\ \nabla^2\phi(x) &= \sum_{i=1}^m \frac{1}{f_i^2(x)} \nabla f_i(x) \nabla f_i(x)^T + \sum_{i=1}^m \frac{1}{-f_i(x)} \nabla^2 f_i(x) \end{aligned}$$

- This means that **the objective fn. of (†) is convex and twice differentiable.**

## Building Block: Newton's Method

- Consider unconstrained minimization of convex  $f : \mathbf{R}^n \rightarrow \mathbf{R}$

$$\min f(x)$$

- Optimality condition:  $\nabla f(x^*) = 0$ .

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### *Newton's Method*

**given** an initial point  $x \in \text{dom} f$ , a tolerance  $\epsilon > 0$ .

**repeat**

1. *Compute the Newton step.*

$$\Delta x = -(\nabla^2 f(x))^{-1} \nabla f(x)$$

2. *Line search.* Choose step size  $\alpha$  by solving

$$\min_{0 \leq \alpha \leq 1} f(x + \alpha \Delta x)$$

3. *Update.*  $x := x + \alpha \Delta x$ .

4. *Stopping criterion.* **quit** if  $\frac{1}{2} \nabla f(x)^T (\nabla^2 f(x))^{-1} \nabla f(x) \leq \epsilon$
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- The idea: first-order approximation of the optimality condition

$$\nabla f(x + \Delta x) \approx \nabla f(x) + \nabla^2 f(x) \Delta x = 0$$

- Newton's method can be extended to handle

$$\begin{aligned} \min \quad & f(x) \\ \text{s.t.} \quad & Ax = b \end{aligned}$$

(Intuitively this is reasonable because the equality constrained min. problem is equivalent to  $\min_z f(A^\dagger b + Fz)$  for some  $F$  such that  $\mathcal{R}(F) = \mathcal{N}(A)$ ).

- Newton's method has fast convergence in general.



# Central Path

- The log barrier approximation

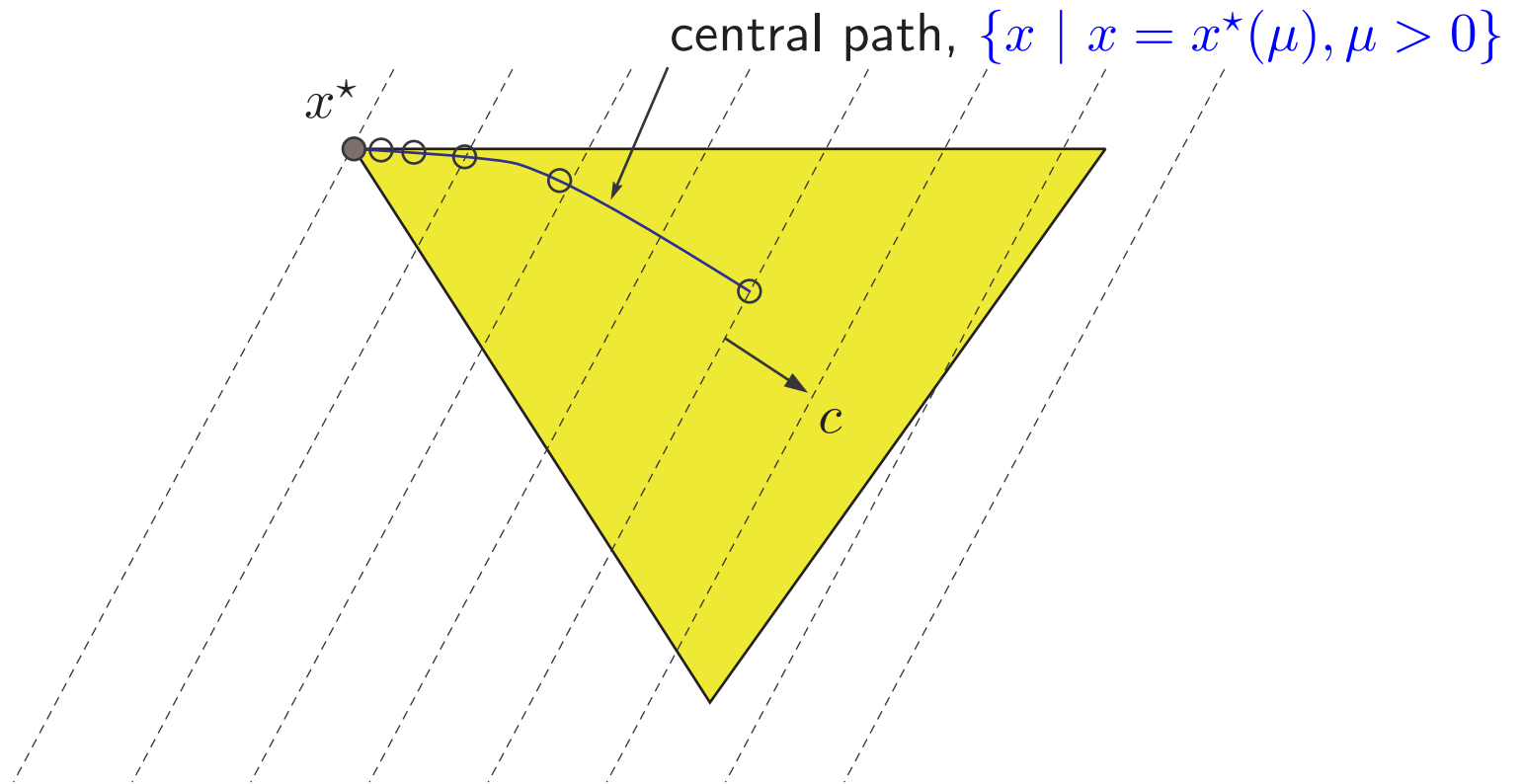
$$\begin{aligned} \min \quad & f_0(x) + \mu\phi(x) \\ \text{s.t.} \quad & Ax = b \end{aligned} \tag{*}$$

- is an accurate approx. of the original problem for  $\mu \rightarrow 0$  (with  $\mu > 0$ ).
- becomes difficult to minimize as  $\mu \rightarrow 0$  (true at least for Newton's method)
- **The basic idea:** start with a large  $\mu$ , and iteratively reduce  $\mu$  until a desired solution accuracy is reached.
- Define  $x^*(\mu)$  to be the solution of (\*).
- **Central path**

$$\{x \mid x = x^*(\mu), \mu > 0\}$$

is the collection of optimal points for various  $\mu$ . A property:

$$f_0(x^*(\mu)) - p^* \leq m\mu$$



**Illustration of the central path of LP. An interior-point method would follow the central path to iteratively approach the optimal solution.**

# A Path Following Method (or Barrier Method)

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## *Path Following Method*

**given** an initial strictly feasible point  $x$ ,  $\mu$ ,  $\epsilon > 0$ , &  $\beta < 1$ .

**repeat**

1. *Centering step.* Starting at  $x$ , use Newton's method to solve

$$x^*(\mu) = \min_x \mu f_0(x) + \phi(x) \text{ s.t. } Ax = b,$$

2. *Update.*  $x := x^*(\mu)$ .

3. *Stopping criterion.* **quit** if  $m\mu < \epsilon$ .

4. *Target shifting.*  $\mu := \beta\mu$ .

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- **Short-step path following:** choose  $\beta$  close to 1.
  - small number of Newton steps per outer iteration
  - large number of outer iterations
- **Long-step path following:** choose a small  $\beta$ .
  - increased number of Newton steps per outer iteration
  - smaller number of outer iterations

# Log Barrier for Conic Optimization

- Conic problem in standard form

$$\begin{aligned} \min \quad & c^T x \\ \text{s.t.} \quad & Ax = b, \quad x \succeq_K 0 \end{aligned}$$

- Log barriers for

- **LP**:  $K = \mathbf{R}_+^n$

$$\phi(x) = -\sum_{i=1}^n \log(x_i)$$

- **SOCP**:  $K = \{(x, t) \mid \|x\|_2 \leq t\}$

$$\phi(x) = \log(t - \|x\|_2^2/t)$$

- **SDP**:  $K = \{X \in \mathbf{S}^n \mid X \succeq 0\}$

$$\phi(X) = -\log \det(X)$$

- These barriers are **self-concordant** (def. skipped), from which appealing convergence results can be proven.

# Primal-Dual Path Following Method for SDP

- Primal-dual path following has good control over solution accuracy.
- **The idea:** Approximate the KKT conditions of the barrier problems.
- Primal-dual path following can be applied to general convex problems, but here we use SDP as an example.
- Features of primal-dual path following
  - simultaneously produce primal and dual points at each iteration
  - the central path parameter  $\mu$  is adaptively updated according to the current primal and dual points.
  - no inner-outer iterations as in primal path following

- SDP in standard form

$$\begin{aligned} \min_X \quad & \mathbf{tr}(CX) \\ \text{s.t.} \quad & X \succeq 0, \mathbf{tr}(A_i X) = b_i, \quad i = 1, \dots, m \end{aligned} \quad (\text{SDP})$$

where  $A_i \in \mathbf{S}^n$ ,  $b_i \in \mathbf{R}$ .

- The dual problem

$$\begin{aligned} \max_{y, Z} \quad & b^T y \\ \text{s.t.} \quad & Z \succeq 0, \quad C - Z - \sum_{i=1}^m y_i A_i = 0 \end{aligned} \quad (\text{DSDP})$$

- The KKT conditions: for  $(X, y, Z) = (X^*, y^*, Z^*)$ ,

$$\begin{aligned} X &\succeq 0, Z \succeq 0 \\ C - Z - \sum_{i=1}^m y_i A_i &= 0 \\ b_i - \mathbf{tr}(A_i X) &= 0, \quad i = 1, \dots, m \\ ZX &= 0 \quad (\text{complementary slackness}) \end{aligned}$$

- Consider the barrier problem of the dual SDP (DSDP)

$$\begin{aligned} \max_{y, Z \succ 0} \quad & b^T y + \mu \log \det Z \\ \text{s.t.} \quad & C - Z - \sum_{i=1}^m y_i A_i = 0 \end{aligned} \tag{DBP}$$

- The KKT conditions for (DBP): for  $(X, y, Z) = (X^*(\mu), y^*(\mu), Z^*(\mu))$ ,

$$\begin{aligned} \mu Z^{-1} - X &= 0 \\ C - Z - \sum_{i=1}^m y_i A_i &= 0 \\ b_i - \text{tr}(A_i X) &= 0, \quad i = 1, \dots, m \end{aligned}$$

- A property for the duality gap:

$$\text{tr}(C X^*(\mu)) - b^T y^*(\mu) = \text{tr}(Z^*(\mu) X(\mu)) = \mu/n$$

As  $\mu \rightarrow 0$ , the duality gap is zero and  $(X^*(\mu), y^*(\mu), Z^*(\mu)) \rightarrow (X^*, y^*, Z^*)$ .

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### Primal-Dual Path Following Method for SDP

**given** an initial strictly feasible point  $(X, y, Z) = (X^{(0)}, y^{(0)}, Z^{(0)})$ ,  $\epsilon > 0$ .

**repeat**

1. *Compute the current  $\mu$ .*  $\mu := \text{tr}(ZX)/n$ .
2. *Target shifting.*  $\mu := \mu/2$ .
3. *Compute search directions for new  $\mu$ .* Compute  $(\Delta X, \Delta y, \Delta Z)$  by solving the KKT eqns.

$$\mu(Z + \Delta Z)^{-1} - (X + \Delta X) = 0$$

$$C - (Z + \Delta Z) - \sum_{i=1}^m (y_i + \Delta y_i) A_i = 0$$

$$b_i - \text{tr}(A_i(X + \Delta X)) = 0, \quad i = 1, \dots, m$$

in an approximate manner.

4. *Line search.* Compute primal & dual step-sizes  $\alpha_p, \alpha_d$  such that

$$X + \alpha_p \Delta X \succ 0 \quad Z + \alpha_d \Delta Z \succ 0$$

5. *Update.*  $X := X + \alpha_p \Delta X$ ,  $y := y + \alpha_d \Delta y$ ,  $Z := Z + \alpha_d \Delta Z$ .

6. *Stopping criterion.* **quit** if  $\text{tr}(ZX) < \epsilon$ .



### Approximation of the KKT equations (Step 3):

For a primal-dual feasible iterate  $(X, y, Z)$ , the eqns in Step 3 are expressed as

$$\mu(Z + \Delta Z)^{-1} = (X + \Delta X) \quad (1)$$

$$-\Delta Z - \sum_{i=1}^m \Delta y_i A_i = 0 \quad (2)$$

$$\text{tr}(A_i \Delta X) = 0, \quad i = 1, \dots, m \quad (3)$$

(2)-(3) are linear & are easy to handle. The difficulty mainly lies in (1).

Apply a 1st-order approximation to (1) (note: from this point there are many possible variations); e.g., in **[Helmberg et. al'96]**,

$$\begin{aligned} \mu I &= (Z + \Delta Z)(X + \Delta X) = ZX + Z\Delta X + X\Delta Z + \Delta Z\Delta X \\ &\approx ZX + Z\Delta X + X\Delta Z \end{aligned} \quad (4)$$

Equations (2)-(4) (modified linearized KKT eqns.) are linear in  $(\Delta X, \Delta y, \Delta Z)$ , and can be solved in closed form.

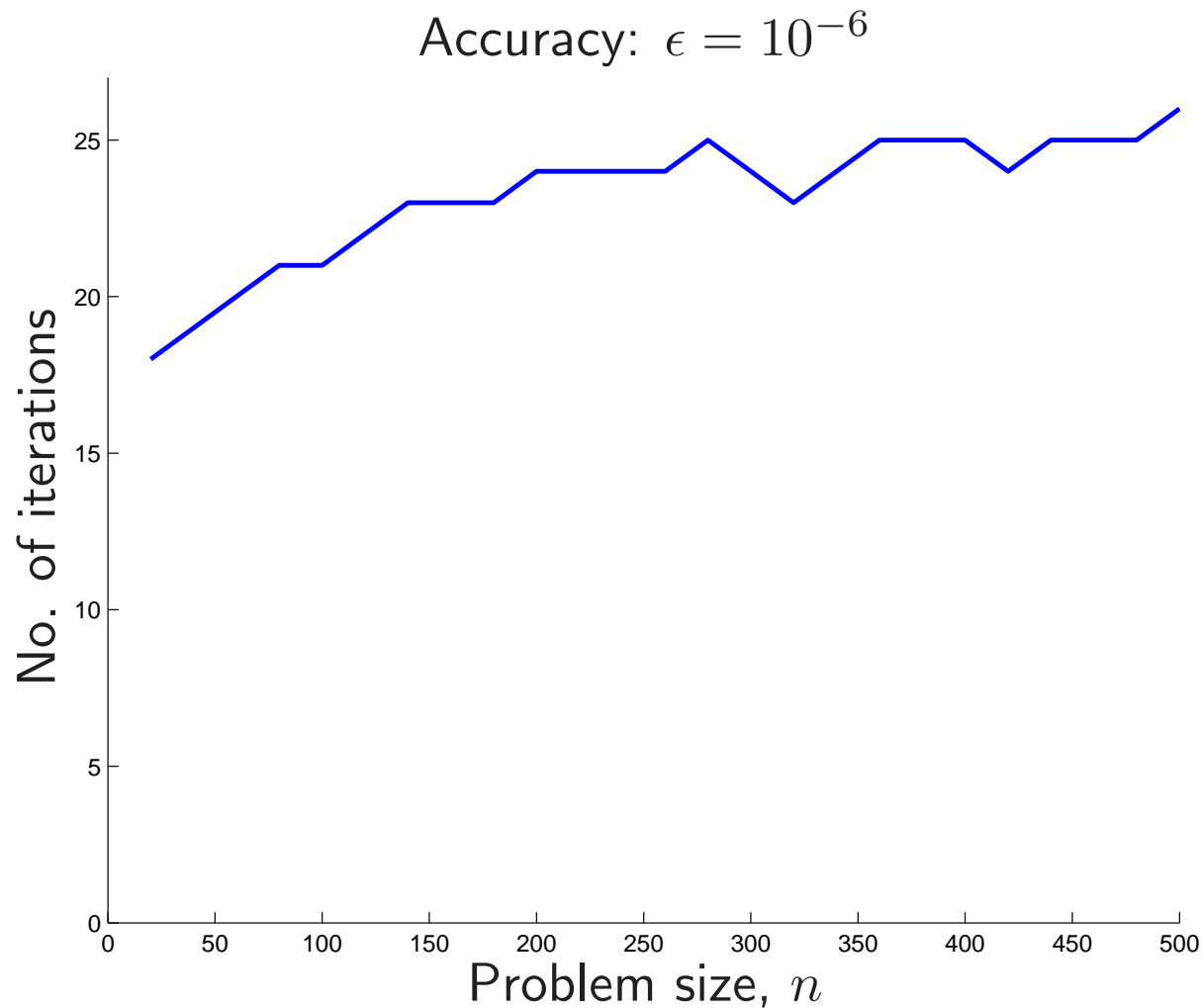
- Convergence analysis for primal-dual path following methods suggests the following:

For a given tolerance  $\epsilon > 0$  (such that  $\text{tr}(CX) - p^* < \epsilon$ ), terminates in

$$\mathcal{O} \left( \sqrt{n} \log \frac{\text{tr}(X^{(0)} Z^{(0)})}{\epsilon} \right) \text{ iterations}$$

in the worst case.

- Practical experience (e.g., with the MIMO detection application described later) suggests that the no. of iterations grows much more slower than  $\sqrt{n}$ .



Nos. of iterations of the primal-dual SDP path following method in practice. The problem is

$$\min \operatorname{tr}(CX) \text{ s.t. } X_{ii} = 1, i = 1, \dots, n$$

As seen, the nos. of iterations appear to be flat for large  $n$ .

## Other (Even More Advanced) Interior-Point Methods

- Self-dual embedding (used in SeDuMi)
- Dual scaling (used in DSDP)
- Infeasible primal-dual path following (used in SDPT3)
- ...

## References

**[Boyd-Vandenberghe'04]** Stephen Boyd and Lieven Vandenberghe, *Convex Optimization*, Cambridge Univ. Press, 2004.

**[Helmberg-Rendl-Vanderbei-Wolkowicz'96]** C. Helmberg, F. Rendl, R.J. Vanderbei, and H. Wolkowicz, "An interior-point method for SDP", *SIAM J. Optim.*, 1996.